

# A REMARK ON PARITY SEQUENCES

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For an integer  $n \geq 2$ , let  $T_n$  be the unique set of positive integers such that:

- (1)  $1 \in T_n$ ;
- (2) if  $t > 1$ , then  $t \in T_n$  iff exactly one of  $t-1$ ,  $t-n$  is in  $T_n$ .

Condition (2) can be rephrased as

**The Triple Criterion:** If  $t \neq 1$ , then  $|\{t-n, t-1, t\} \cap T_n| \in \{0, 2\}$ .

If  $n=2$ , then the set  $T_n$  is closely related to the Fibonacci sequence; specifically,  $t \in T_2$  iff the  $t^{\text{th}}$  term of the Fibonacci sequence is odd.

We ask, for each  $n$ , which numbers are uniquely expressible as the sum of two distinct elements of  $T_n$ . In general, for any given  $n$ , one can determine exactly which numbers are uniquely expressible. If  $n=2$ , it is easy to see that there are five such numbers:  $3=1+2$ ,  $5=1+4$ ,  $7=2+5$ ,  $8=1+7$ , and  $10=2+8$ . If  $n=3$ , then there are exactly eight uniquely expressible numbers:  $3=1+2$ ,  $4=1+3$ ,  $5=2+3$ ,  $6=1+5$ ,  $7=2+5$ ,  $8=3+5$ ,  $9=1+8$ , and  $16=1+15$ . If  $n=4$ , then there are exactly five uniquely expressible numbers:  $3=1+2$ ,  $4=1+3$ ,  $6=2+4$ ,  $8=2+6$ , and  $16=4+12$ . If  $n \geq 3$ , then  $1, 2, 3 \in T_n$ , so that 3 and 4 are uniquely expressible.

The principal theorem of this note answers this question for all other situations. Let  $U_n$  be the set of all integers which are uniquely expressible as the sum of two distinct elements of  $T_n$ . Thus, we have just observed that

$$U_2 = \{3, 5, 7, 8, 10\}, U_3 = \{3, 4, 5, 6, 7, 8, 9, 16\}, \text{ and } U_4 = \{3, 4, 6, 8, 16\}.$$

The following principal theorem characterizes  $U_n$  for  $n \geq 5$ .

**Theorem:** Let  $n \geq 5$ . Then  $U_n = \{3, 4, n^2 - n + 3, 2n^2 - 2n + 4\}$  if  $n = 2^k + 1$  for some  $k$ , and  $U_n = \{3, 4\}$  otherwise.

The remainder of this paper consists of two sections. The first contains a discussion of the motivation for the principal theorem, and the second contains its proof. The second section can be read independently of the first.

## 1. MOTIVATION

For an integer  $n \geq 2$ , let  $f_1, f_2, f_3, \dots$  be the sequence defined by the initial conditions

$$f_1 = f_2 = \dots = f_n = 1$$

and the recurrence relation

$$f_{n+j} = f_j + f_{n+j-1}$$

for  $j \geq 1$ . If, in particular,  $n=2$ , then the Fibonacci sequence has just been defined, and, as another example, if  $n=5$ , then we get the sequence

$$1, 1, 1, 1, 1, 2, 3, 4, 5, 6, 8, 11, 15, 20, 26, 34, 45, 60, 80, 106, \dots$$

From this sequence, we define another sequence  $t_1, t_2, t_3, \dots$ , which we will call the  $n^{\text{th}}$  parity sequence: we set  $t_i = j$  iff the  $i^{\text{th}}$  odd term in the sequence  $f_1, f_2, f_3, \dots$  is  $f_j$ . For example, the 5<sup>th</sup> parity sequence is

$$1, 2, 3, 4, 5, 7, 9, 12, 13, 17, 22, 23, 24, \dots$$

Then  $T_n = \{t_1, t_2, t_3, \dots\}$ .

The principal theorem extends the result of [4] but in a somewhat disguised form. What is essentially proved in [4] is this theorem weakened by requiring that  $n$  be an even number, thereby eliminating any exceptional cases.

We next discuss some background for the result of [4] and, consequently, of the above theorem. For positive integers  $u < v$ , the 1-additive sequence based on  $u, v$  is the sequence  $s_1, s_2, s_3, \dots$ , where  $s_1 = u$ ,  $s_2 = v$ , and  $s_{n+2}$  is the least  $a > s_{n+1}$  for which there is a unique pair of integers  $i, j$  such that  $1 \leq i < j \leq n+1$  and  $a = s_i + s_j$ . For example, the 1-additive sequence based on 1, 2 is the sequence

$$1, 2, 3, 4, 6, 8, 11, 13, 16, 18, 26, 28, \dots,$$

which was introduced by Ulam [5]. This sequence is still not well understood, but it appears to have a quite erratic behavior. Other 1-additive sequences, such as the one based on 2, 3 also exhibit a similar erratic behavior. In contrast to this, the 1-additive sequence based on 2,  $v$ , where  $v \geq 5$  is an odd number, has a much more predictable behavior.

Finch made the definition in [2] that the (increasing) sequence  $s_1, s_2, s_3, \dots$  is *regular* if there are positive integers  $m, p$ , and  $d$  such that whenever  $i \geq m$ , then  $s_{i+p} = s_i + d$ . (He refers to the least such  $p$  as the *period* of the sequence and to the least such  $d$  as the *fundamental difference*.) He observed in [2] that a 1-additive sequence having only finitely many even terms is regular. He then went on to make the conjecture, based on extensive numerical evidence, that for relatively prime  $u < v$ , the 1-additive sequence based on  $u, v$  has only finitely many even terms iff one of the following holds:

- (i)  $u = 2$  and  $v \geq 5$  is odd;
- (ii)  $u = 4$  and  $v \geq 5$  is odd;
- (iii)  $u = 5$  and  $v = 6$ ;
- (iv)  $u \geq 6$  is even;
- (v)  $u \geq 7$  is odd and  $v$  is even.

For each of the cases (i)-(v), he made a conjecture as to what the finite sets are. For example, in (i) the set of even terms is  $\{2, 2v+2\}$ , and in (ii) the set is  $\{4, 2v+4, 4v+4\}$  provided that  $v \neq 2^m - 1$  for any  $m \geq 3$ . The conjecture for (i) was proved correct in [4], and for (ii) it was proved correct in [1] in the case  $v \equiv 1 \pmod{4}$ . For (iii) the set is

$$\{6, 16, 26, 36, 80, 124, 144, 172, 184, 196, 238, 416, 448\},$$

and in this case the truth of the conjecture can be verified by direct computation.

Now suppose that  $D = \{d_1, d_2, \dots, d_k\}$  is a finite set of integers, where  $d_1 < d_2 < \dots < d_k$ . Let us say for now that the sequence  $t_1, t_2, t_3, \dots$  is the 1-incremental sequence based on  $D$  if  $t_1 = 1$  and  $t_{n+1}$  is the least  $a > t_n$  for which there is a unique pair of integers  $i, j$  such that  $1 \leq i \leq n$ ,  $1 \leq j \leq k$ , and  $a = t_i + d_j$ . For example, the 1-incremental sequence based on  $\{1, 5\}$  is

$$1, 2, 3, 4, 5, 7, 9, 12, 13, 17, 22, 23, 24, \dots$$

Notice that this sequence is identical to the 5<sup>th</sup> parity sequence. In general, the  $n^{\text{th}}$  parity sequence is identical to the 1-incremental sequence based on  $\{1, n\}$ .

The connection between 1-incremental sequences and the regularity of 1-additive sequences, elaborating on Finch's observation [2], will be discussed next.

Consider the 1-additive sequence  $s_1, s_2, s_3, \dots$  based on  $u, v$ , where  $u = 2d_1$  is even and  $v$  is odd. Suppose that  $2d_1, 2d_2, \dots, 2d_k$  are all the even terms that are no greater than  $2(d_{k-1} + d_k)$  occurring in the 1-additive sequence, where  $d_1 < d_2 < \dots < d_k$ . Let  $t_1, t_2, t_3, \dots$  be the 1-incremental sequence based on  $D = \{d_1, d_2, \dots, d_k\}$  and let  $T = \{t_1, t_2, t_3, \dots\}$ . It is easy to check that

$$\{s_1, s_2, s_3, \dots\} = \{2t + v - 2 : t \in T\} \cup \{2d_1, 2d_2, \dots, 2d_k\}.$$

Now consider 1-additive sequences based on  $2, v$ , where  $v \geq 5$  is an odd integer. The result of [4] is thus seen to be equivalent to the principal theorem restricted to even  $n \geq 6$ . This leads naturally to the question that this theorem answers.

Every  $n^{\text{th}}$  parity sequence is regular. (In fact, it is obvious that every 1-incremental sequence is regular.) However, even a little more is true for these sequences (and for all 1-incremental sequences based on 2-element sets, as well). Let  $P(n)$  be the period of the  $n^{\text{th}}$  parity sequence  $t_1, t_2, t_3, \dots$ , and let  $D(n)$  be the fundamental difference. Then, it follows from the Triple Criterion that, for each  $i \geq 1$ ,  $t_{i+P(n)} = t_i + D(n)$ . Also  $D(n)$  is the least  $d > 1$  for which none of  $d, d-1, d-2, \dots, d-n+2$  is in  $T_n$ . Tabulation of  $2D(n)$  and  $P(n)$  for many even  $n \geq 6$  can be found in [3].

## 2. THE PROOF

We will need an analysis of the  $(2^k + 1)^{\text{th}}$  parity sequence. An analysis of the  $(2^k)^{\text{th}}$  parity sequence was given in [4]. As a comparison, we summarize that analysis here.

**Proposition 1 ([4]):** Let  $k \geq 1$  and let  $n = 2^k$ . Let  $1 \leq t \leq 4n^2$  and suppose that  $t = 2in + j$ , where  $0 \leq i < 2n$  and  $1 \leq j \leq 2n$ . Then:

- (1) if  $i < n$  and  $j \leq n$ , then  $t \in T_{2n}$  iff  $in + j \in T_n$ ;
- (2) if  $i < n$  and  $j > n$ , then  $t \in T_{2n}$  iff  $in + j - n \in T_n$ ;
- (3) if  $i \geq n$  and  $j \leq n$ , then  $t \in T_{2n}$  iff  $(i - n)n + j \in T_n$  and  $j < n$ ;
- (4) if  $i \geq n$  and  $j > n$ , then  $t \in T_{2n}$  iff  $j = 2n$ .  $\square$

The following notation from Section 1 will be used. Recall from Section 1 that, for each  $n \geq 2$ , there is  $d \geq 1$  such that, for any  $t \geq 1$ ,  $t \in T_n$  iff  $t + d \in T_n$ . We let  $D(n)$  be the least such  $d$ . Clearly,  $D(n)$  is the least  $d \geq 1$  such that  $d + 1, d + 2, d + 3, \dots, d + n \in T_n$ , and also it is the least  $d \geq 1$  such that  $d, d - 1, d - 2, \dots, d - (n - 2) \notin T_n$ .

Using Proposition 1, we can easily prove by induction that, if  $n = 2^k$ , then the following hold: if  $1 \leq i \leq n$ , then  $in \in T_n$ ; if  $1 \leq j \leq n$ , then  $(n - 1)j \in T_n$ ; if  $i < n$  and  $n - i \leq j < n$ , then  $in + j \notin T_n$ . From this it follows that  $n^2 - 1$  is the least  $d \geq 1$  such that  $\{d, d - 1, d - 2, \dots, d - n + 2\} \cap T_n = \emptyset$ . Thus,  $D(n) = 4^k - 1 = n^2 - 1$ . It can also be shown that  $P(n) = 3^k - 1$ .

There is another way to characterize the elements of  $T_{2^k}$ . We introduce some notation. For nonnegative integers  $t$  and  $i$ , we let  $b_i(t)$  be the  $i^{\text{th}}$  digit in the binary expansion of  $t$ . For example, since  $37 = 1 + 4 + 32$ , we get that  $b_i(37) = 1$  if  $i = 0, 2, 5$  and  $b_i(37) = 0$  for all other nonnegative integers  $i$ .

**Proposition 2:** Suppose  $k \geq 1$  and  $n = 2^k$ , and let  $1 \leq t \leq n^2 = 2^{2k}$ . Then  $t \in T_n$  iff whenever  $0 \leq r < k$ , then  $b_r(t) \cdot b_{k+r}(t) = 0$ .

*Proof:* Let us first consider the special case of the proposition when  $b_{k-1}(t) = 1$ ,  $b_{2k-1}(t) = 0$ , and  $b_r(t) = 0$  for all  $r < k-1$ . Clearly,  $b_r(t) \cdot b_{k+r}(t) = 0$  for all  $r < k$ . It is easily checked by induction on  $k$  that Proposition 1 implies that all such  $t$  are in  $T_n$ .

We now turn to the proof of the proposition in general. The proof is by induction on  $k$ . For  $k = 1$ , it is easily checked. Let  $n = 2^k$ ; we will prove it for the case  $2n = 2^{k+1}$ . Let  $1 \leq t \leq 4n^2$ , and (as in Proposition 1) let  $t = 2in + j$ , where  $0 \leq i < 2n$  and  $1 \leq j \leq 2n$ . The proof splits naturally into the same four cases as does Proposition 1. Since each one is routine, we will do just case (1), where  $i < n$  and  $j \leq n$ . Notice that these restrictions on  $i$  and  $j$  are equivalent to the condition that  $b_k(t-1) = b_{2k+1}(t-1) = 0$ , and this condition splits into two subcases.

**Subcase 1:**  $b_k(t) = b_{2k+1}(t) = 0$  and  $b_r(t) = 1$  for some  $r < k$ . Since  $b_k(t) = 0$ , we need only be concerned with  $b_r(t) \cdot b_{(k+1)+r}(t)$  for  $r < k$ . For such  $r$ ,  $b_r(t) = b_r(in + j)$  and  $b_{(k+1)+r}(t) = b_{k+r}(in + j)$ , so the result easily follows from the inductive hypothesis.

**Subcase 2:**  $b_k(t) = 1$ ,  $b_{2k+1}(t) = 0$ , and  $b_r(t) = 0$  for all  $r < k$ . But this is just the special case that was noted at the beginning of the proof.  $\square$

In ways analogous to those in Propositions 1 and 2, the sets  $T_{2^{k+1}}$  can be analyzed. This is done in Propositions 3 and 4, respectively.

**Proposition 3:** Let  $k \geq 0$  and let  $n = 2^k$ . Let  $1 \leq t \leq (2n+1)^2$  and suppose that  $t = i(2n+1) + j$ , where  $0 \leq i \leq 2n$  and  $1 \leq j \leq 2n+1$ . Then:

- (1) if  $i \leq n$  and  $j \leq n+1$ , then  $t \in T_{2n+1}$  iff  $i(n+1) + j \in T_{n+1}$ ;
- (2) if  $i \leq n$  and  $j > n+1$ , then  $t \in T_{2n+1}$  iff  $i(n+1) + j - n \in T_{n+1}$  and  $i \neq n$ ;
- (3) if  $i > n$  and  $j \leq n+1$ , then  $t \in T_{2n+1}$  iff  $(i-n)(n+1) + j \in T_{n+1}$ ;
- (4) if  $i > n$  and  $j > n+1$ , then  $t \in T_{2n+1}$  iff  $i = 2n$ .

*Proof:* The proof is by induction on  $k$ . For  $k = 0$ , it is easily checked. Consider some  $k > 0$ , and assume, as the inductive hypothesis, that the proposition holds for all smaller values of  $k$ . Let  $n = 2^k$ , and let  $t = i(2n+1) + j$ , where  $0 \leq i \leq 2n$  and  $1 \leq j \leq 2n+1$ . We proceed by induction on  $t$ . The proof splits naturally into four cases. Since each is routine, we will show only case (1), where  $i \leq n$  and  $j \leq n+1$ . This case splits into three subcases.

**Subcase 1:**  $i = 0$ . Then  $t = j$ , and it is clear that  $j \in T_{2n+1}$  and  $j \in T_{n+1}$ .

**Subcase 2:**  $i > 0$  and  $j > 1$ . Then, using the Triple Criterion and the inductive hypothesis on  $t$ , we see that  $t \in T_{2n+1}$  iff

$$\begin{aligned}
 & t-1 \in T_{2n+1} \Leftrightarrow t-(2n+1) \notin T_{2n+1} \\
 \text{iff} & \\
 & i(2n+1) + j - 1 \in T_{2n+1} \Leftrightarrow (i-1)(2n+1) + j \notin T_{2n+1} \\
 \text{iff} & \\
 & i(n+1) + j - 1 \in T_{n+1} \Leftrightarrow (i-1)(n+1) + j \notin T_{n+1} \\
 \text{iff} & \\
 & i(n+1) + j \in T_{n+1}.
 \end{aligned}$$

**Subcase 3:  $i > 0$  and  $j = 1$ .** Then, again using the Triple Criterion and the inductive hypothesis on  $t$ , we see that  $t \in T_{2n+1}$  iff

$$t - 1 \in T_{2n+1} \Leftrightarrow t - (2n + 1) \notin T_{2n+1}$$

iff

$$(i - 1)(2n + 1) + (2n + 1) \in T_{2n+1} \Leftrightarrow (i - 1)(2n + 1) + 1 \notin T_{2n+1}$$

iff

$$i(n + 1) \in T_{n+1} \Leftrightarrow (i - 1)(n + 1) + 1 \notin T_{n+1}$$

iff

$$i(n + 1) + 1 \in T_{n+1}. \quad \square$$

**Proposition 4:** Suppose  $k \geq 1$  and  $n = 2^k$ , and let  $2 \leq t \leq n^2 + 1$ . Then  $t \in T_{n+1}$  iff whenever  $0 \leq r < k$ , then  $b_r(t - 2) \geq b_{k+r}(t - 2)$ .

**Proof:** The proof is by induction on  $k$ . For small values of  $k$ , say  $k = 1, 2$ , it is easily checked. Let  $n = 2^k$ ; we will prove it for the case  $2n = 2^{k+1}$ . Let  $2 \leq t \leq 4n^2 + 1$ , and (as in Proposition 3) let  $t = i(2n + 1) + j$ , where  $0 \leq i \leq 2n$  and  $1 \leq j \leq 2n + 1$ . As  $t \geq 2$ , it is obvious that  $2 \leq i + j$ . The proof splits naturally into the same four cases as does Proposition 3. Since each one is routine, we will show just case (1), where  $i \leq n$  and  $j \leq n + 1$ . Thus,  $2 \leq i + j \leq 2n + 1 = 2^{k+1} + 1$ .

**Subcase 1:  $i + j < 2^k$ .** Since  $t = i2^{k+1} + (i + j)$ , where  $2 \leq i + j < 2^k$ , it is clear that  $b_k(t - 2) = b_{2k+1}(t - 2) = 0$  and also that  $b_r(t - 2) = b_r(in + (i + j) - 2)$  and  $b_{k+(r+1)}(t - 2) = b_{k+r}(in + (i + j) - 2)$  for  $r < k$ . Therefore, from the inductive hypothesis,

$$t \in T_{2n+1} \Leftrightarrow i(n + 1) + j \in T_{n+1} \Leftrightarrow b_r(i(n + 1) + j - 2) \geq b_{k+r}(i(n + 1) + j - 2) \\ \text{for } r < k \Leftrightarrow b_r(t - 2) \geq b_{(k+1)+r}(t - 2) \text{ for } r \leq k.$$

**Subcase 2:  $i + j = 2^k$ .** Then  $b_0(t - 2) = b_k(t - 2) = b_{2k+1}(t - 2) = 0$ , and  $b_r(t - 2) = 1$  if  $1 \leq r < k$ . Also,  $b_{k+1}(t - 2) = 0$  iff  $i$  is even. Therefore, we have that  $b_r(t - 2) \geq b_{(k+1)+r}(t - 2)$  whenever  $0 \leq r \leq k$  iff  $i$  is even. On the other hand,

$$t \in T_{2n+1} \Leftrightarrow i(n + 1) + j \in T_{n+1} \Leftrightarrow (i + 1)n \in T_{n+1} \Leftrightarrow b_k((i + 1)n - 2) = 0 \Leftrightarrow i \text{ is even.}$$

**Subcase 3:  $i + j = 2^k + 1$ .** Then  $b_k(t - 2) = b_{2k+1}(t - 2) = 0$  and  $b_r(t - 2) = 1$  if  $0 \leq r < k$ . Thus, we have that  $b_r(t - 2) \geq b_{(k+1)+r}(t - 2)$  whenever  $0 \leq r \leq k$ . On the other hand,

$$i2^{k+1} + 2^k + 1 \in T_{2n+1} \Leftrightarrow (i + 1)n + 1 \in T_{n+1},$$

which is the case since  $b_r((i + 1)n - 1) = 1$  for all  $r < k$ .

**Subcase 4:  $2^k + 2 \leq i + j < 2^{k+1}$ .** As in Subcase 1, it is clear that  $b_k(t - 2) = 1$  and also that  $b_r(t - 2) = b_r(in + (i + j) - 2)$  and  $b_{k+(r+1)}(t - 2) = b_{k+r}(in + (i + j) - 2)$  for  $r < k$ . Therefore, from the inductive hypothesis,

$$t \in T_{2n+1} \Leftrightarrow i(n + 1) + j \in T_{n+1} \Leftrightarrow b_r(i(n + 1) + j - 2) \geq b_{k+r}(i(n + 1) + j - 2) \\ \text{for } r < k \Leftrightarrow b_r(t - 2) \geq b_{(k+1)+r}(t - 2) \text{ for } r \leq k.$$

**Subcase 5:  $i + j = 2^{k+1}$ .** (This subcase is similar to Subcase 2.) Then  $b_0(t - 2) = b_k(t - 2) = 0$  and  $b_r(t - 2) = 1$  if  $1 \leq r < k$ . Also,  $b_{k+1}(t - 2) = 0$  iff  $i$  is even. Therefore, we have that  $b_r(t - 2) \geq b_{(k+1)+r}(t - 2)$  whenever  $0 \leq r \leq k$  iff  $i$  is even. On the other hand,

$$t \in T_{2n+1} \Leftrightarrow i(n+1) + j \in T_{n+1} \Leftrightarrow (i+1)n \in T_{n+1} \Leftrightarrow b_k((i+1)n-2) = 0 \Leftrightarrow i \text{ is even.}$$

**Subcase 6:**  $i + j = 2^{k+1} + 1$ . Therefore, we have  $i = n$ ,  $j = n + 1$ , and  $t = 2^{2k+1} + 2^{k+1} + 1$ . Then  $b_r(t-2) = 1$  for all  $r \leq k$ . Thus, we have that  $b_r(t-2) \geq b_{(k+1)+r}(t-2)$  whenever  $0 \leq r \leq k$ . On the other hand,

$$t \in T_{2n+1} \Leftrightarrow n(2n+1) + (n+1) \in T_{2n+1} \Leftrightarrow n(n+1) + (n+1)T_{n+1} \Leftrightarrow 2^{2k} + 2^{k+1} + 1 \in T_{n+1},$$

which is the case by the inductive hypothesis since  $b_r(2^{2k} + 2^{k+1} - 1) = 1$  for all  $r < k$ .  $\square$

**Proposition 5:** Suppose that  $n \geq 2$  and  $s = in + j$ , where  $0 \leq i < n$  and  $0 \leq j < n$ . Then:

- (1) if  $i < n - 1$  and  $j < n - i - 1$ , then  $D(n) - s \notin T_n$ ;
- (2) if  $i < n$  and  $j = n - i - 1$ , then  $D(n) - s \in T_n$ ;
- (3) if  $i < n - 1$  and  $j = n - 1$ , then  $D(n) - s \in T_n$ .

**Proof:** The proof is by induction on  $s$ . We provide the details. We let  $s = in + j$ , where  $0 \leq i < n$  and either  $0 \leq j < n - i - 1$  or  $j = n - 1$ . Suppose the proposition is true for all smaller values of  $s$ . Let  $a = D(n) - s$ , so  $a$  might be negative. We will determine whether or not  $a \in T_n$  by seeing whether or not each of  $a + n$  and  $a + n - 1$  is in  $T_n$ , and then use the Triple Criterion applied to  $\{a, a + n - 1, a + n\}$ . To do so, it is necessary to know that  $a + n \neq 1$ . In each case, it will be clear that  $a + n \neq 1$  since there will be  $b$  such that  $a < b < a + n$  and  $b \notin T_n$ .

**Case 1:**  $i = 0$ ,  $0 \leq j < n - 1$ . Then  $a + n = n + D(n) - j \in T_n$  since  $n - j \in T_n$ , and  $a + n - 1 = n + D(n) - j - 1 \in T_n$  since  $n - j - 1 \in T_n$ . Therefore,  $a \notin T_n$ .

**Case 2:**  $i = 0$ ,  $j = n - 1$ . Then  $a + n = D(n) + 1 \in T_n$  since  $1 \in T_n$ , and  $a + n - 1 = D(n) \notin T_n$  by the inductive hypothesis. Therefore,  $a \in T_n$ .

**Case 3:**  $0 < i < n - 1$ ,  $j = 0$ . Then  $a + n = D(n) = (i - 1)n \notin T_n$  and  $a + n - 1 = D(n) - ((i - 1)n + 1) \notin T_n$  by the inductive hypothesis. Therefore,  $a \notin T_n$ .

**Case 4:**  $i = n - 1$ ,  $j = 0$ . Then  $a + n = D(n) - (n - 2)n \notin T_n$  and  $a + n - 1 = D(n) - ((n - 2)n + 1) \in T_n$  by the inductive hypothesis. Therefore,  $a \in T_n$ .

**Case 5:**  $0 < i < n - 1$ ,  $0 < j < n - i - 1$ . Then  $a + n = D(n) - ((i - 1)n + j) \notin T_n$  and  $a + n - 1 = D(n) - ((i - 1)n + (j + 1)) \notin T_n$  by the inductive hypothesis. Therefore,  $a \notin T_n$ .

**Case 6:**  $0 < i < n - 1$ ,  $j = n - i - 1$ . Then  $a + n = D(n) - ((i - 1)n + j) \notin T_n$  and  $a + n - 1 = D(n) - ((i - 1)n + (j + 1)) \in T_n$  by the inductive hypothesis. Therefore,  $a \in T_n$ .

**Case 7:**  $0 < i < n - 1$ ,  $j = n - 1$ . Then  $a + n = D(n) - ((i - 1)n + (n - 1)) \in T_n$  and  $a + n - 1 = D(n) - in \notin T_n$  by the inductive hypothesis. Therefore,  $a \in T_n$ .  $\square$

Two special instances of Proposition 5 will be used later on. If  $i = 1$ , then (2) shows that  $D(n) - 2n + 2 \in T_n$  and (3) shows that  $D(n) - 2n + 1 \in T_n$ .

**Corollary 6:** Let  $n \geq 2$ .

- (1) Then  $D(n) \geq n^2 - n + 1$ .
- (2) If  $n = 2^k + 1$ , then  $D(n) = n^2 - n + 1$ .

**Proof:** It follows from Proposition 5(2) (letting  $i = n - 1$ ,  $j = 0$ ) that  $D(n) - (n - 1)n \in T_n$ , so that  $D(n) \geq n^2 - n + 1$ . For  $n = 2^k + 1$ , it follows from Proposition 4 that, if  $n^2 - n + 2 \leq t \leq n^2 + 1$ , then  $t \in T_n$ , so that  $D(n) \leq n^2 - n + 1$ .  $\square$

It can be shown that, if  $n = 2^k + 1$ , then  $P(n) = 3^k + 1$ .

It follows that, if  $n = 2^k + 1$ , then  $n^2 - n + 2 = 1 + D(n) \in T_n$  and  $2n^2 - 2n + 3 = 1 + 2D(n) \in T_n$ . We can now deduce a part of the principal theorem.

**Corollary 7:** Suppose  $k \geq 1$  and  $n = 2^k + 1$ . Let  $a, b \in T_n$  be such that  $a < b$ .

(1) If  $a + b = n^2 - n + 3$ , then  $a = 1$  and  $b = n^2 - n + 2$ .

(2) If  $a + b = 2n^2 - 2n + 4$ , then  $a = 1$  and  $b = 2n^2 - 2n + 3$ .

**Proof:** Let  $a, b \in T_n$  such that  $a < b$ .

(1) Suppose  $a + b = n^2 - n + 3$  but  $a > 1$ . Let  $c = a - 2$ ,  $d = b - 2$ , and  $e = c + d = n^2 - n - 1$ . Then  $b_{n-1}(e) = 1$  and, for  $0 \leq i < 2^k = n - 1$ ,  $b_i(e) = 1$  iff  $i < 2^{k-1}$ . Since  $b \in T_n$  and  $b \leq n^2 - n + 1$ , it must be that  $b \leq n^2 - 2n + 1$ , so that  $d \leq n^2 - 2n$ . Therefore, there is  $j < k$  such that  $b_{k+j}(c) = 1$  and then, also,  $b_j(d) = 1$ . Consider some such  $j$ . Clearly, for each  $i < k$ ,  $b_i(c) \neq b_i(d)$ . It is also clear that, if  $k \leq i < 2k$ , then  $b_i(c) = b_i(d)$ . But then  $1 = b_{k+j}(c) = b_{k+j}(d) = b_j(d) \neq b_j(c)$ , contradicting Proposition 4.

(2) Suppose  $a + b = 2n^2 - 2n + 4$ , but  $a > 1$ . Then  $b \geq n^2 - n + 3$ . Let  $c = b - (n^2 - n + 1)$ , so that  $c \geq 2$ ,  $c \in T_n$  by Corollary 6(2), and  $a + c = n^2 - n + 3$ . It follows from (1) that  $a = c$ , which is impossible because  $a + c$  is odd.  $\square$

With the assistance of Proposition 3 or Proposition 4 we can, in general, determine a large initial segment of any  $n^{\text{th}}$  parity sequence.

**Proposition 8:** Let  $k \geq 0$ ,  $q \geq 1$ ,  $n = q2^k + 1$ , and  $m = 2^k + 1$ . Suppose  $1 \leq t \leq n(m - 1)$ , and let  $t = in + j$ , where  $0 \leq i < m - 1$  and  $1 \leq j \leq n$ . Let  $j = r2^k + s$ , where  $0 \leq r < q$  and  $1 \leq s \leq m$ . Then  $t \in T_n$  iff  $im + s \in T_m$ .

**Proof:** The proof is a straightforward induction on  $t$ .  $\square$

**Proof of the Theorem:** Suppose that  $n \geq 5$ . As previously observed,  $3, 4 \in U_n$ . It follows from Corollary 7 that, if  $n = 2^k + 1$ , then  $n^2 - n + 3$  and  $2n^2 - 2n + 4$  are in  $U_n$ .

For the reverse inclusion, suppose that  $a, b \in T_n$  are such that  $a < b$ ,  $a + b \geq 5$ , and for no  $a', b' \in T_n$  is it the case that  $a \neq a' < b' \neq b$  and  $a' + b' = a + b$ .

We can assume that  $a + b > 2n$ . (For, as is easy to check, if  $s \leq 2n$ , then the number of pairs  $a, b \in T_n$  such that  $a < b$  and  $a + b = s$  is  $\lfloor \frac{1}{2}(s - 1) \rfloor$  if  $s \leq n$ , is  $\frac{1}{2}(n - 1)$  if  $s > n$  is odd, is  $\frac{n}{2}$  if  $s > n$  and  $n$  is even, and is  $\frac{1}{2}(n - 2)$  if  $s > n$  is even and  $n$  is even.) Since  $\{1, 2, 3, \dots, n\} \subseteq T_n$  and since  $\{a + b - 1, a + b - 2, a + b - 3, \dots, a + b - n\} \cap T_n \neq \emptyset$ , it must be that  $1 \leq a \leq n$ . Also,  $b \leq 2D(n) + n$ , as other-wise setting  $a' = a + D(n)$  and  $b' = b - D(n)$  yields a contradiction.

Now let  $n = q2^k + 1$ , where  $q$  is odd. We consider two cases.

**$a = 1$ :** Then  $\{b - 1, b - 2, b - 3, \dots, b - n + 1\} \cap T_n = \emptyset$ . Thus,  $b = 1 + pD(n)$  for some  $p \geq 1$ , and also  $p \leq 2$ , as otherwise  $a' = 1 + D(n)$ ,  $b' = 1 + (p - 1)D(n)$  would yield a contradiction. By Corollary 7, we can suppose that  $q > 1$ . Then, from Proposition 8, we get that  $2^k + 1 \in T_n$  and, from Proposition 5, that  $D(n) - 2^k + 1 \in T_n$ . It follows from Corollary 6 that  $D(n) - 2^k + 1 > 2^k + 1$ . Thus, setting  $a' = 2^k + 1$  and  $b' = b - 2^k + 1$  yields a contradiction.

$1 < a \leq n$ : Then  $\{a+b-i: 1 \leq i \leq n \text{ and } i \neq a\} \cap T_n = \emptyset$ . Thus, Proposition 5 implies that  $b = pD(n) - n + 1$  for some  $p > 0$ . Either  $a+n \in T_n$  or  $a+n-1 \in T_n$ . Let  $a'$  be whichever one is in  $T_n$ , and let  $b' = b - (a' - a)$ . Then, by Proposition 5,  $b' \in T_n$ , thereby arriving at a contradiction.  $\square$

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